

Section 6.3

Formal Reasoning

- A *formal proof* (or derivation) is a sequence of wffs, where each wff is either a premise or the result of applying a proof rule to certain previous wffs in the sequence.

Basic Proof Rules

$\frac{A, B}{A \wedge B}$ Conjunction (Conj)

$\frac{A \wedge B}{A} \frac{A \wedge B}{B}$ Simplification (Simp)

$\frac{A}{A \vee B} \frac{B}{A \vee B}$ Addition (Add)

$\frac{A \vee B, \neg A}{B} \frac{A \vee B, \neg B}{A}$ Disjunctive Syllogism (DS)

More Basic Proof Rules

$\frac{A \rightarrow B, A}{B}$ Modus Ponens (MP)

$\frac{\text{From } A, \text{ derive } B}{A \rightarrow B}$ Conditional Proof (CP)

$\frac{\neg\neg A}{A} \frac{A}{\neg\neg A}$ Double Negation (DN)

$\frac{A, \neg A}{\text{False}}$ Contradiction (Contr)

$\frac{\text{From } \neg A, \text{ derive False}}{A}$ Indirect Proof (IP)

First example proof

- Put each wff on a numbered line along with a reason. Use the letter P for a premise and follow the proof with QED.
- Example: Prove that the argument with the premises $A \vee C \rightarrow D$, $\neg B$, $A \vee B$ and conclusion D is valid.

$$1. A \vee C \rightarrow D \quad P$$

First example proof

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- Example: Prove that the argument with the premises $A \vee C \rightarrow D$, $\neg B$, $A \vee B$ and conclusion D is valid.

1. $A \vee C \rightarrow D$ P
2. $\neg B$ P

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- Example: Prove that the argument with the premises $A \vee C \rightarrow D$, $\neg B$, $A \vee B$ and conclusion D is valid.

1. $A \vee C \rightarrow D$ P
2. $\neg B$ P
3. $A \vee B$ P

First example proof

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- Example: Prove that the argument with the premises $A \vee C \rightarrow D$, $\neg B$, $A \vee B$ and conclusion D is valid.

1.	$A \vee C \rightarrow D$	P
2.	$\neg B$	P
3.	$A \vee B$	P
4.	A	2,3,DS

First example proof

- Put each wff on a numbered line along with a reason. Use the letter P for a premise and follow the proof with QED.
- Example: Prove that the argument with the premises $A \vee C \rightarrow D$, $\neg B$, $A \vee B$ and conclusion D is valid.

1.	$A \vee C \rightarrow D$	P
2.	$\neg B$	P
3.	$A \vee B$	P
4.	A	2,3,DS
5.	$A \vee C$	4, Add

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- Put each wff on a numbered line along with a reason. Use the letter P for a premise and follow the proof with QED.
- Example: Prove that the argument with the premises $A \vee C \rightarrow D$, $\neg B$, $A \vee B$ and conclusion D is valid.

1.	$A \vee C \rightarrow D$	P
2.	$\neg B$	P
3.	$A \vee B$	P
4.	A	2,3,DS
5.	$A \vee C$	4, Add
6.	D	1,5,MP

First example proof

- Put each wff on a numbered line along with a reason. Use the letter P for a premise and follow the proof with QED.
- Example: Prove that the argument with the premises $A \vee C \rightarrow D$, $\neg B$, $A \vee B$ and conclusion D is valid.

1.	$A \vee C \rightarrow D$	P
2.	$\neg B$	P
3.	$A \vee B$	P
4.	A	2,3,DS
5.	$A \vee C$	4, Add
6.	D	1,5,MP
	QED	

Using CP

- If a proof consists of a derivation from a premise A to a conclusion B that does not contain any uses of CP or IP, then we can apply CP to obtain a tautology $A \rightarrow B$. The reason we obtain a tautology is that the proof rules used in the derivation are valid arguments. So the truth of A implies the truth of B , which tells us that $A \rightarrow B$ is a tautology.
- When using CP in this way, instead of writing $A \rightarrow B$, we'll write QED along with the line numbers of the derivation followed by CP.

CP proof

Example: Prove that $(A \vee C \rightarrow D) \wedge \neg B \wedge (A \vee B) \rightarrow D$ is a tautology.

$$1. A \vee C \rightarrow D \quad P$$

CP proof

Example: Prove that $(A \vee C \rightarrow D) \wedge \neg B \wedge (A \vee B) \rightarrow D$ is a tautology.

1. $A \vee C \rightarrow D$ P
2. $\neg B$ P

CP proof

Example: Prove that $(A \vee C \rightarrow D) \wedge \neg B \wedge (A \vee B) \rightarrow D$ is a tautology.

1. $A \vee C \rightarrow D$ P
2. $\neg B$ P
3. $A \vee B$ P

CP proof

Example: Prove that $(A \vee C \rightarrow D) \wedge \neg B \wedge (A \vee B) \rightarrow D$ is a tautology.

1. $A \vee C \rightarrow D$ P
2. $\neg B$ P
3. $A \vee B$ P
4. A 2,3,DS

CP proof

Example: Prove that $(A \vee C \rightarrow D) \wedge \neg B \wedge (A \vee B) \rightarrow D$ is a tautology.

1. $A \vee C \rightarrow D$ P
2. $\neg B$ P
3. $A \vee B$ P
4. A 2,3,DS
5. $A \vee C$ 4, Add

CP proof

Example: Prove that $(A \vee C \rightarrow D) \wedge \neg B \wedge (A \vee B) \rightarrow D$ is a tautology.

1. $A \vee C \rightarrow D$ P
2. $\neg B$ P
3. $A \vee B$ P
4. A 2,3,DS
5. $A \vee C$ 4, Add
6. D 1,5,MP

CP proof

Example: Prove that $(A \vee C \rightarrow D) \wedge \neg B \wedge (A \vee B) \rightarrow D$ is a tautology.

- | | | |
|----|--------------------------|---------|
| 1. | $A \vee C \rightarrow D$ | P |
| 2. | $\neg B$ | P |
| 3. | $A \vee B$ | P |
| 4. | A | 2,3,DS |
| 5. | $A \vee C$ | 4, Add |
| 6. | D | 1,5,MP |
| | QED | 1-6,CP. |

Another CP proof

Example: Prove $(A \vee B \rightarrow C \wedge D) \wedge A \wedge (C \rightarrow E) \rightarrow D \wedge E$ is a tautology.

1. $A \vee B \rightarrow C \wedge D$ P

Another CP proof

Example: Prove $(A \vee B \rightarrow C \wedge D) \wedge A \wedge (C \rightarrow E) \rightarrow D \wedge E$ is a tautology.

1. $A \vee B \rightarrow C \wedge D$ P
2. A P

Another CP proof

Example: Prove $(A \vee B \rightarrow C \wedge D) \wedge A \wedge (C \rightarrow E) \rightarrow D \wedge E$ is a tautology.

1. $A \vee B \rightarrow C \wedge D$ P
2. A P
3. $C \rightarrow E$ P

Another CP proof

Example: Prove $(A \vee B \rightarrow C \wedge D) \wedge A \wedge (C \rightarrow E) \rightarrow D \wedge E$ is a tautology.

1. $A \vee B \rightarrow C \wedge D$ P
2. A P
3. $C \rightarrow E$ P
4. $A \vee B$ 2,Add

Another CP proof

Example: Prove $(A \vee B \rightarrow C \wedge D) \wedge A \wedge (C \rightarrow E) \rightarrow D \wedge E$ is a tautology.

1. $A \vee B \rightarrow C \wedge D$ P
2. A P
3. $C \rightarrow E$ P
4. $A \vee B$ 2, Add
5. $C \wedge D$ 1,4, MP

Another CP proof

Example: Prove $(A \vee B \rightarrow C \wedge D) \wedge A \wedge (C \rightarrow E) \rightarrow D \wedge E$ is a tautology.

- | | | |
|----|-----------------------------------|---------|
| 1. | $A \vee B \rightarrow C \wedge D$ | P |
| 2. | A | P |
| 3. | $C \rightarrow E$ | P |
| 4. | $A \vee B$ | 2, Add |
| 5. | $C \wedge D$ | 1,4, MP |
| 6. | C | 5, Simp |

Another CP proof

Example: Prove $(A \vee B \rightarrow C \wedge D) \wedge A \wedge (C \rightarrow E) \rightarrow D \wedge E$ is a tautology.

- | | | |
|----|-----------------------------------|---------|
| 1. | $A \vee B \rightarrow C \wedge D$ | P |
| 2. | A | P |
| 3. | $C \rightarrow E$ | P |
| 4. | $A \vee B$ | 2, Add |
| 5. | $C \wedge D$ | 1,4, MP |
| 6. | C | 5, Simp |
| 7. | E | 3,6, MP |

Another CP proof

Example: Prove $(A \vee B \rightarrow C \wedge D) \wedge A \wedge (C \rightarrow E) \rightarrow D \wedge E$ is a tautology.

1. $A \vee B \rightarrow C \wedge D$ P
2. A P
3. $C \rightarrow E$ P
4. $A \vee B$ 2, Add
5. $C \wedge D$ 1,4, MP
6. C 5, Simp
7. E 3,6, MP
8. D 5, Simp

Another CP proof

Example: Prove $(A \vee B \rightarrow C \wedge D) \wedge A \wedge (C \rightarrow E) \rightarrow D \wedge E$ is a tautology.

- | | | |
|----|-----------------------------------|-----------|
| 1. | $A \vee B \rightarrow C \wedge D$ | P |
| 2. | A | P |
| 3. | $C \rightarrow E$ | P |
| 4. | $A \vee B$ | 2, Add |
| 5. | $C \wedge D$ | 1,4, MP |
| 6. | C | 5, Simp |
| 7. | E | 3,6, MP |
| 8. | D | 5, Simp |
| 9. | $D \wedge E$ | 7,8, Conj |

Another CP proof

Example: Prove $(A \vee B \rightarrow C \wedge D) \wedge A \wedge (C \rightarrow E) \rightarrow D \wedge E$ is a tautology.

- | | | |
|----|-----------------------------------|-----------|
| 1. | $A \vee B \rightarrow C \wedge D$ | P |
| 2. | A | P |
| 3. | $C \rightarrow E$ | P |
| 4. | $A \vee B$ | 2, Add |
| 5. | $C \wedge D$ | 1,4, MP |
| 6. | C | 5, Simp |
| 7. | E | 3,6, MP |
| 8. | D | 5, Simp |
| 9. | $D \wedge E$ | 7,8, Conj |
| | QED | 1-9, CP. |

Subproofs

- A *subproof* is a proof that is part of another proof. It always starts with a new premise and always ends by applying CP or IP to the derivation from that premise. When this happens, the premise is *discharged* and the wffs from the derivation become inactive.
- Indent the statements of the subproof and write down the result of CP or IP without indentation.
- Example: We'll prove that $(A \vee B) \rightarrow (\neg B \rightarrow A)$ is a tautology.

1.	$A \vee B$	P
2.	$\neg B$	P [for $\neg B \rightarrow A$]
3.	A	1,2,DS
4.	$\neg B \rightarrow A$	2,3, CP
	QED	1,4,CP.

Another subproof example

Prove: $(A \vee B \rightarrow C) \rightarrow (A \rightarrow C) \wedge (B \rightarrow C)$ is a tautology.

1. $A \vee B \rightarrow C$ P

Another subproof example

Prove: $(A \vee B \rightarrow C) \rightarrow (A \rightarrow C) \wedge (B \rightarrow C)$ is a tautology.

- | | | |
|----|--------------------------|----------------------------|
| 1. | $A \vee B \rightarrow C$ | P |
| 2. | A | P [for $A \rightarrow C$] |

Another subproof example

Prove: $(A \vee B \rightarrow C) \rightarrow (A \rightarrow C) \wedge (B \rightarrow C)$ is a tautology.

- | | | |
|----|--------------------------|----------------------------|
| 1. | $A \vee B \rightarrow C$ | P |
| 2. | A | P [for $A \rightarrow C$] |
| 3. | $A \vee B$ | 2, Add |

Another subproof example

Prove: $(A \vee B \rightarrow C) \rightarrow (A \rightarrow C) \wedge (B \rightarrow C)$ is a tautology.

- | | | |
|----|--------------------------|----------------------------|
| 1. | $A \vee B \rightarrow C$ | P |
| 2. | A | P [for $A \rightarrow C$] |
| 3. | $A \vee B$ | 2, Add |
| 4. | C | 1,3, MP |

Another subproof example

Prove: $(A \vee B \rightarrow C) \rightarrow (A \rightarrow C) \wedge (B \rightarrow C)$ is a tautology.

1.	$A \vee B \rightarrow C$	P
2.	A	P [for $A \rightarrow C$]
3.	$A \vee B$	2, Add
4.	C	1,3, MP
5.	$A \rightarrow C$	2-4, CP

Another subproof example

Prove: $(A \vee B \rightarrow C) \rightarrow (A \rightarrow C) \wedge (B \rightarrow C)$ is a tautology.

1.	$A \vee B \rightarrow C$	P
2.	A	P [for $A \rightarrow C$]
3.	$A \vee B$	2,Add
4.	C	1,3,MP
5.	$A \rightarrow C$	2-4, CP
6.	B	P [for $B \rightarrow C$]

Another subproof example

Prove: $(A \vee B \rightarrow C) \rightarrow (A \rightarrow C) \wedge (B \rightarrow C)$ is a tautology.

1.	$A \vee B \rightarrow C$	P
2.	A	P [for $A \rightarrow C$]
3.	$A \vee B$	2, Add
4.	C	1,3, MP
5.	$A \rightarrow C$	2-4, CP
6.	B	P [for $B \rightarrow C$]
7.	$A \vee B$	6, Add

Another subproof example

Prove: $(A \vee B \rightarrow C) \rightarrow (A \rightarrow C) \wedge (B \rightarrow C)$ is a tautology.

1.	$A \vee B \rightarrow C$	P
2.	A	P [for $A \rightarrow C$]
3.	$A \vee B$	2, Add
4.	C	1,3, MP
5.	$A \rightarrow C$	2-4, CP
6.	B	P [for $B \rightarrow C$]
7.	$A \vee B$	6, Add
8.	C	1,7, MP

Another subproof example

Prove: $(A \vee B \rightarrow C) \rightarrow (A \rightarrow C) \wedge (B \rightarrow C)$ is a tautology.

1.	$A \vee B \rightarrow C$	P
2.	A	P [for $A \rightarrow C$]
3.	$A \vee B$	2, Add
4.	C	1,3, MP
5.	$A \rightarrow C$	2-4, CP
6.	B	P [for $B \rightarrow C$]
7.	$A \vee B$	6, Add
8.	C	1,7, MP
9.	$B \rightarrow C$	6-8, CP

Another subproof example

Prove: $(A \vee B \rightarrow C) \rightarrow (A \rightarrow C) \wedge (B \rightarrow C)$ is a tautology.

1.	$A \vee B \rightarrow C$	P
2.	A	P [for $A \rightarrow C$]
3.	$A \vee B$	2,Add
4.	C	1,3,MP
5.	$A \rightarrow C$	2-4, CP
6.	B	P [for $B \rightarrow C$]
7.	$A \vee B$	6, Add
8.	C	1,7,MP
9.	$B \rightarrow C$	6-8, CP
10.	$(A \rightarrow C) \wedge (B \rightarrow C)$	5,9,Conj

Another subproof example

Prove: $(A \vee B \rightarrow C) \rightarrow (A \rightarrow C) \wedge (B \rightarrow C)$ is a tautology.

1.	$A \vee B \rightarrow C$	P
2.	A	P [for $A \rightarrow C$]
3.	$A \vee B$	2,Add
4.	C	1,3,MP
5.	$A \rightarrow C$	2-4, CP
6.	B	P [for $B \rightarrow C$]
7.	$A \vee B$	6, Add
8.	C	1,7,MP
9.	$B \rightarrow C$	6-8, CP
10.	$(A \rightarrow C) \wedge (B \rightarrow C)$	5,9,Conj
	QED	1,5,9-10,CP.

An Indirect Proof

- If a proof consists of a derivation from a premise $\neg A$ to the conclusion False, then we could apply CP to obtain $\neg A \rightarrow \text{False}$. But we also know that $A \equiv \neg A \rightarrow \text{False}$. So the result of the derivation is A. This is the IP rule.
- Example: We'll prove the tautology $\neg(A \wedge \neg A)$.
 1. $\neg\neg(A \wedge \neg A)$ P [for $\neg(A \wedge \neg A)$]

An Indirect Proof

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- Example: We'll prove the tautology $\neg(A \wedge \neg A)$.
 1. $\neg\neg(A \wedge \neg A)$ P [for $\neg(A \wedge \neg A)$]
 2. $A \wedge \neg A$ 1,DN

An Indirect Proof

- If a proof consists of a derivation from a premise $\neg A$ to the conclusion False, then we could apply CP to obtain $\neg A \rightarrow \text{False}$. But we also know that $A \equiv \neg A \rightarrow \text{False}$. So the result of the derivation is A. This is the IP rule.
- Example: We'll prove the tautology $\neg(A \wedge \neg A)$.
 1. $\neg\neg(A \wedge \neg A)$ P [for $\neg(A \wedge \neg A)$]
 2. $A \wedge \neg A$ 1, DN
 3. A 2, Simp

An Indirect Proof

- If a proof consists of a derivation from a premise $\neg A$ to the conclusion False, then we could apply CP to obtain $\neg A \rightarrow \text{False}$. But we also know that $A \equiv \neg A \rightarrow \text{False}$. So the result of the derivation is A. This is the IP rule.
- Example: We'll prove the tautology $\neg(A \wedge \neg A)$.
 1. $\neg\neg(A \wedge \neg A)$ P [for $\neg(A \wedge \neg A)$]
 2. $A \wedge \neg A$ 1, DN
 3. A 2, Simp
 4. $\neg A$ 2, Simp

An Indirect Proof

- If a proof consists of a derivation from a premise $\neg A$ to the conclusion False, then we could apply CP to obtain $\neg A \rightarrow \text{False}$. But we also know that $A \equiv \neg A \rightarrow \text{False}$. So the result of the derivation is A. This is the IP rule.
- Example: We'll prove the tautology $\neg(A \wedge \neg A)$.

1.	$\neg\neg(A \wedge \neg A)$	P [for $\neg(A \wedge \neg A)$]
2.	$A \wedge \neg A$	1, DN
3.	A	2, Simp
4.	$\neg A$	2, Simp
5.	False	3,4, Contr

An Indirect Proof

- If a proof consists of a derivation from a premise $\neg A$ to the conclusion False, then we could apply CP to obtain $\neg A \rightarrow \text{False}$. But we also know that $A \equiv \neg A \rightarrow \text{False}$. So the result of the derivation is A. This is the IP rule.
- Example: We'll prove the tautology $\neg(A \wedge \neg A)$.

1.	$\neg\neg(A \wedge \neg A)$	P [for $\neg(A \wedge \neg A)$]
2.	$A \wedge \neg A$	1, DN
3.	A	2, Simp
4.	$\neg A$	2, Simp
5.	False	3,4, Contr
	QED	1-5, IP.

IP proofs

- IP is most often used in a subproof setting when proving a conditional of the form $V \rightarrow W$. Start with V as a premise for a CP proof. Then start an IP subproof with premise $\neg W$. When a contradiction is reached, we obtain W by IP. Then CP gives the result $V \rightarrow W$.
- As with CP subproofs, the result of IP is written with no indentation.

IP Example

Prove the tautology: $(A \rightarrow B) \wedge (A \vee B) \rightarrow B$.

1. $A \rightarrow B$ P

IP Example

Prove the tautology: $(A \rightarrow B) \wedge (A \vee B) \rightarrow B$.

1. $A \rightarrow B$ P
2. $A \vee B$ P

IP Example

Prove the tautology: $(A \rightarrow B) \wedge (A \vee B) \rightarrow B$.

1. $A \rightarrow B$ P
2. $A \vee B$ P
3. $\neg B$ P [for B]

IP Example

Prove the tautology: $(A \rightarrow B) \wedge (A \vee B) \rightarrow B$.

1. $A \rightarrow B$ P
2. $A \vee B$ P
3. $\neg B$ P [for B]
4. A 2,3,DS

IP Example

Prove the tautology: $(A \rightarrow B) \wedge (A \vee B) \rightarrow B$.

1. $A \rightarrow B$ P
2. $A \vee B$ P
3. $\neg B$ P [for B]
4. A 2,3,DS
5. B 1,4,MP

IP Example

Prove the tautology: $(A \rightarrow B) \wedge (A \vee B) \rightarrow B$.

1. $A \rightarrow B$ P
2. $A \vee B$ P
3. $\neg B$ P [for B]
4. A 2,3,DS
5. B 1,4,MP
6. False 3,5,Contr

IP Example

Prove the tautology: $(A \rightarrow B) \wedge (A \vee B) \rightarrow B$.

1. $A \rightarrow B$ P
2. $A \vee B$ P
3. $\neg B$ P [for B]
4. A 2,3,DS
5. B 1,4,MP
6. False 3,5,Contr
7. B 3-6, IP

IP Example

Prove the tautology: $(A \rightarrow B) \wedge (A \vee B) \rightarrow B$.

1. $A \rightarrow B$ P
 2. $A \vee B$ P
 3. $\neg B$ P [for B]
 4. A 2,3,DS
 5. B 1,4,MP
 6. False 3,5,Contr
 7. B 3-6, IP
- QED 1-2,7,CP.

Another IP Example

We'll prove the converse of $(A \vee B) \rightarrow (\neg B \rightarrow A)$. Proof of $(\neg B \rightarrow A) \rightarrow (A \vee B)$.

$$1. \neg B \rightarrow A \quad P$$

Another IP Example

We'll prove the converse of $(A \vee B) \rightarrow (\neg B \rightarrow A)$. Proof of $(\neg B \rightarrow A) \rightarrow (A \vee B)$.

1. $\neg B \rightarrow A$ P
2. $\neg(A \vee B)$ P [for $A \vee B$]

Another IP Example

We'll prove the converse of $(A \vee B) \rightarrow (\neg B \rightarrow A)$. Proof of $(\neg B \rightarrow A) \rightarrow (A \vee B)$.

1. $\neg B \rightarrow A$ P
2. $\neg(A \vee B)$ P [for $A \vee B$]
3. $\neg B$ P [for B]

Another IP Example

We'll prove the converse of $(A \vee B) \rightarrow (\neg B \rightarrow A)$. Proof of $(\neg B \rightarrow A) \rightarrow (A \vee B)$.

1. $\neg B \rightarrow A$ P
2. $\neg(A \vee B)$ P [for $A \vee B$]
3. $\neg B$ P [for B]
4. A 1,3,MP

Another IP Example

We'll prove the converse of $(A \vee B) \rightarrow (\neg B \rightarrow A)$. Proof of $(\neg B \rightarrow A) \rightarrow (A \vee B)$.

1. $\neg B \rightarrow A$ P
2. $\neg(A \vee B)$ P [for $A \vee B$]
3. $\neg B$ P [for B]
4. A 1,3,MP
5. $A \vee B$ 4,Add

Another IP Example

We'll prove the converse of $(A \vee B) \rightarrow (\neg B \rightarrow A)$. Proof of $(\neg B \rightarrow A) \rightarrow (A \vee B)$.

1. $\neg B \rightarrow A$ P
2. $\neg(A \vee B)$ P [for $A \vee B$]
3. $\neg B$ P [for B]
4. A 1,3,MP
5. $A \vee B$ 4,Add
6. False 2,5,Contr

Another IP Example

We'll prove the converse of $(A \vee B) \rightarrow (\neg B \rightarrow A)$. Proof of $(\neg B \rightarrow A) \rightarrow (A \vee B)$.

1. $\neg B \rightarrow A$ P
2. $\neg(A \vee B)$ P [for $A \vee B$]
3. $\neg B$ P [for B]
4. A 1,3,MP
5. $A \vee B$ 4,Add
6. False 2,5,Contr
7. B 3-6,IP

Another IP Example

We'll prove the converse of $(A \vee B) \rightarrow (\neg B \rightarrow A)$. Proof of $(\neg B \rightarrow A) \rightarrow (A \vee B)$.

1. $\neg B \rightarrow A$ P
2. $\neg(A \vee B)$ P [for $A \vee B$]
3. $\neg B$ P [for B]
4. A 1,3,MP
5. $A \vee B$ 4,Add
6. False 2,5,Contr
7. B 3-6,IP
8. $A \vee B$ 7,Add

Another IP Example

We'll prove the converse of $(A \vee B) \rightarrow (\neg B \rightarrow A)$. Proof of $(\neg B \rightarrow A) \rightarrow (A \vee B)$.

1. $\neg B \rightarrow A$ P
2. $\neg(A \vee B)$ P [for $A \vee B$]
3. $\neg B$ P [for B]
4. A 1,3,MP
5. $A \vee B$ 4,Add
6. False 2,5,Contr
7. B 3-6,IP
8. $A \vee B$ 7,Add
9. False 2,8,Contr

Another IP Example

We'll prove the converse of $(A \vee B) \rightarrow (\neg B \rightarrow A)$. Proof of $(\neg B \rightarrow A) \rightarrow (A \vee B)$.

- | | | |
|-----|------------------------|---------------------|
| 1. | $\neg B \rightarrow A$ | P |
| 2. | $\neg(A \vee B)$ | P [for $A \vee B$] |
| 3. | $\neg B$ | P [for B] |
| 4. | A | 1,3,MP |
| 5. | $A \vee B$ | 4,Add |
| 6. | False | 2,5,Contr |
| 7. | B | 3-6,IP |
| 8. | $A \vee B$ | 7,Add |
| 9. | False | 2,8,Contr |
| 10. | $A \vee B$ | 2,7-9,IP |

Another IP Example

We'll prove the converse of $(A \vee B) \rightarrow (\neg B \rightarrow A)$. Proof of $(\neg B \rightarrow A) \rightarrow (A \vee B)$.

1. $\neg B \rightarrow A$ P
 2. $\neg(A \vee B)$ P [for $A \vee B$]
 3. $\neg B$ P [for B]
 4. A 1,3,MP
 5. $A \vee B$ 4,Add
 6. False 2,5,Contr
 7. B 3-6,IP
 8. $A \vee B$ 7,Add
 9. False 2,8,Contr
 10. $A \vee B$ 2,7-9,IP
- QED 1,10,CP.

Derived Proof Rules

$\frac{A \rightarrow B, \neg B}{\neg A}$ Modus Tollens (MT)

$\frac{A \vee B, A \rightarrow C, B \rightarrow C}{C}$ Proof by Cases (Cases)

$\frac{A \rightarrow B, B \rightarrow C}{A \rightarrow C}$ Hypothetical Syllogism (HS)

$\frac{A \vee B, A \rightarrow C, B \rightarrow D}{C \vee D}$ Constructive Dilemma (CD)

Proofs using derived rules

We'll give two proof of the tautology
 $(A \rightarrow C) \wedge (B \rightarrow C) \rightarrow (A \vee B \rightarrow C)$.

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Second proof of the same tautology

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4. $\neg C$ P [for C]

Second proof of the same tautology

- | | | |
|----|-------------------|-----------------------------------|
| 1. | $A \rightarrow C$ | P |
| 2. | $B \rightarrow C$ | P |
| 3. | $A \vee B$ | P [for $A \vee B \rightarrow C$] |
| 4. | $\neg C$ | P [for C] |
| 5. | $\neg A$ | 1,4,MT |

Second proof of the same tautology

- | | | |
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| 1. | $A \rightarrow C$ | P |
| 2. | $B \rightarrow C$ | P |
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| 6. | B | 3,5,DS |

Second proof of the same tautology

1.	$A \rightarrow C$	P
2.	$B \rightarrow C$	P
3.	$A \vee B$	P [for $A \vee B \rightarrow C$]
4.	$\neg C$	P [for C]
5.	$\neg A$	1,4,MT
6.	B	3,5,DS
7.	$\neg B$	2,4,MT

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1.	$A \rightarrow C$	P
2.	$B \rightarrow C$	P
3.	$A \vee B$	P [for $A \vee B \rightarrow C$]
4.	$\neg C$	P [for C]
5.	$\neg A$	1,4,MT
6.	B	3,5,DS
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8.	False	6,7,Contr

Second proof of the same tautology

1.	$A \rightarrow C$	P
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3.	$A \vee B$	P [for $A \vee B \rightarrow C$]
4.	$\neg C$	P [for C]
5.	$\neg A$	1,4,MT
6.	B	3,5,DS
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9.	C	4-8,IP

Second proof of the same tautology

1. $A \rightarrow C$ P
2. $B \rightarrow C$ P
3. $A \vee B$ P [for $A \vee B \rightarrow C$]
4. $\neg C$ P [for C]
5. $\neg A$ 1,4,MT
6. B 3,5,DS
7. $\neg B$ 2,4,MT
8. False 6,7,Contr
9. C 4-8,IP
10. $A \vee B \rightarrow C$ 3,9,CP

Second proof of the same tautology

1.	$A \rightarrow C$	P
2.	$B \rightarrow C$	P
3.	$A \vee B$	P [for $A \vee B \rightarrow C$]
4.	$\neg C$	P [for C]
5.	$\neg A$	1,4,MT
6.	B	3,5,DS
7.	$\neg B$	2,4,MT
8.	False	6,7,Contr
9.	C	4-8,IP
10.	$A \vee B \rightarrow C$	3,9,CP
	QED	1-2,10,CP.